

Four Attacks and a Proof for Telegram

Joint Latvian-Estonian Theory Days 2022 May 06, 2022

Martin R. Albrecht, Lenka Mareková, Kenneth G. Paterson, Igors Stepanovs







Based on a paper to appear at IEEE S&P 2022. More information at https://mtpsym.github.io/

Telegram

Monthly active users in Jan 2022:

According to Statistica 2022.



WhatsApp

 $2000 \cdot 10^6$



 \mathbf{WeChat}

 $1263 \cdot 10^6$



FB Messenger $988 \cdot 10^6$



 $574 \cdot 10^6$



Snapchat

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Telegram

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Collective Information Security in Large-Scale Urban Protests: the Case of Hong Kong

Martin R. Albrecht, Jorge Blasco, Rikke Bjerg Jensen, and Lenka Mareková, *Royal Holloway, University of London*

<u>Predominant</u> in Hong Kong protests.

<u>Perceived more secure</u> than competitors.

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Advantages of **Telegram**:

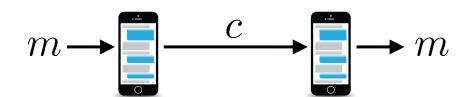
Group chats for up to 200000 people.
Support of <u>pseudonyms</u> in group chats.
Other features:

..., anonymous polls, disappearing messages, timed or scheduled messages, ability to delete messages sent by others, ...



	Cloud Chats	Secret Chats
ENCRYPTION	client-server	end-to-end
GROUPS	✓	X
1-ON-1	\	✓
ENABLED BY DEFAULT		X

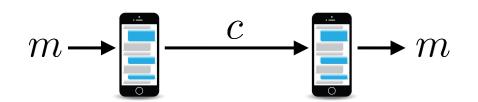
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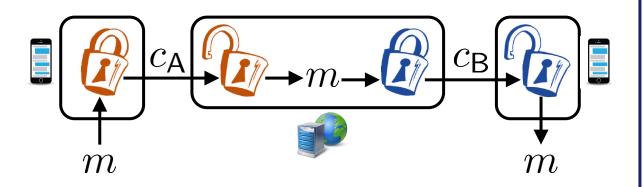
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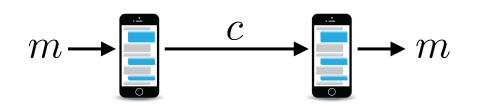






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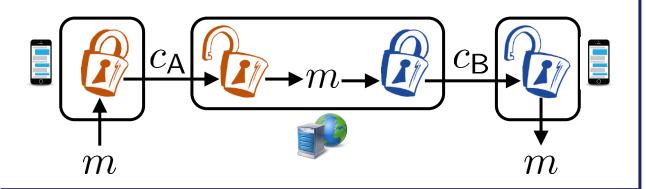
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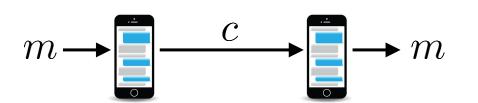


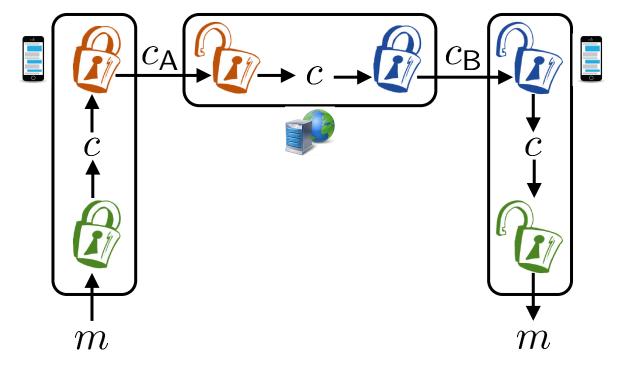


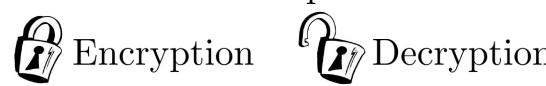


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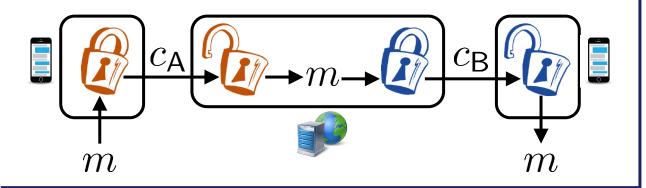
Secret Chats





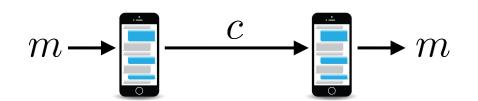


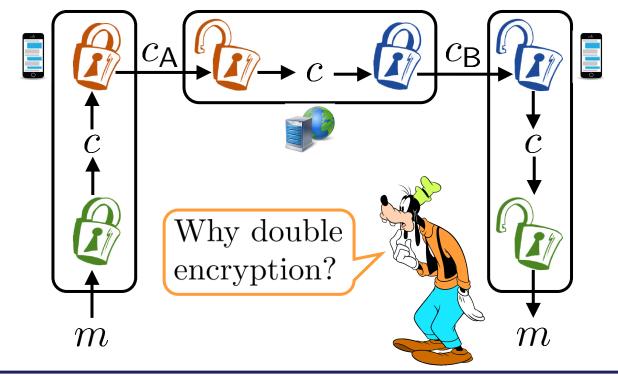


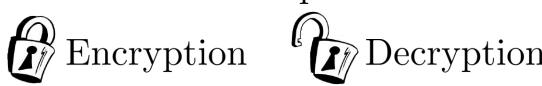


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Secret Chats











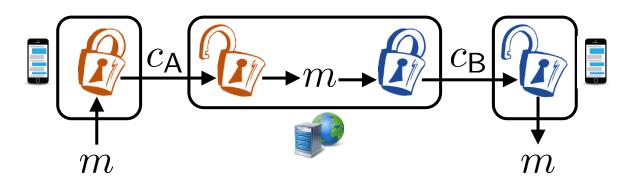
MTProto 2.0

MTProto 2.0 is **Telegram**'s equivalent of the TLS record protocol.

```
Telegram launched with MTProto 1.0.
2016 + MTProto 1.0 is not CCA-secure [JO16].
          Input validation bug (message replay) [SK17]. Telegram released MTProto 2.0.
2018 + Input validation bug in key exchange [K18].
         - MTProto 2.0 secure in symbolic model [MV20].
                (assuming ideal building blocks)
```

MTProto 2.0 is not well-studied.





Why not use TLS?

https://core.telegram.org/techfaq

Telegram FAQ



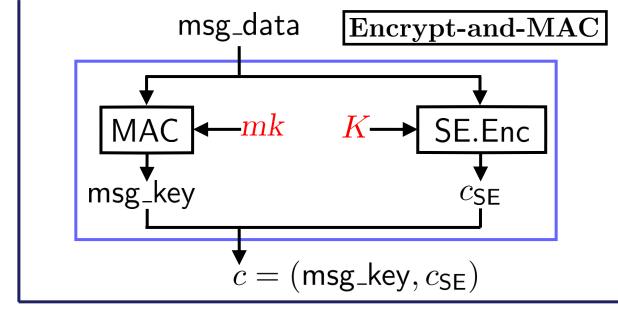
Q: Why are you not using X? (insert solution) While other ways of achieving the same cryptographic goals, undoubtedly, exist, we feel that the present solution is both robust and also succeeds at our secondary task of beating unencrypted messengers in terms of delivery time and stability.

MTProto 2.0

MTProto 2.0 is **Telegram**'s equivalent of the TLS record protocol.

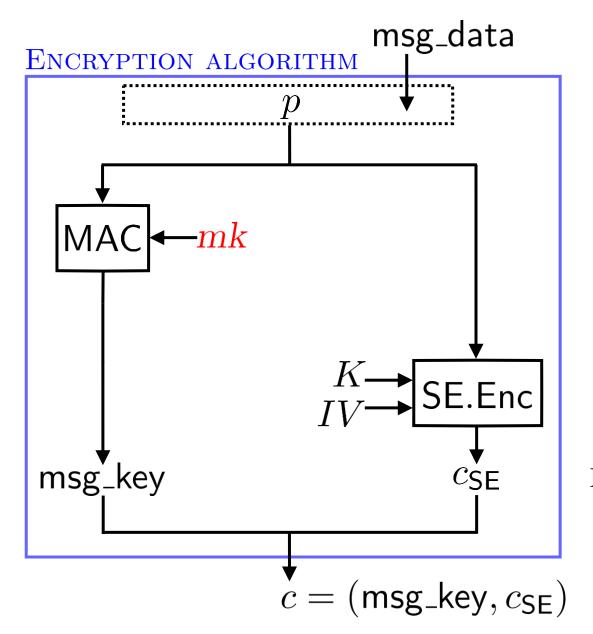
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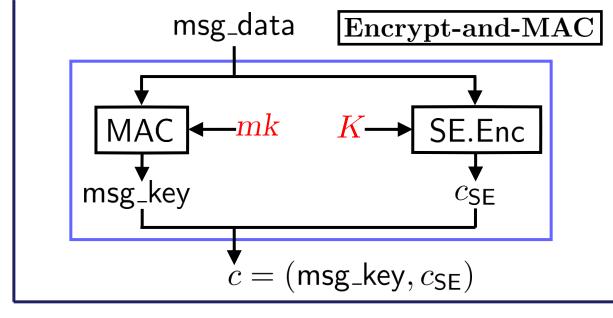
MTProto 2.0 is not well-studied.



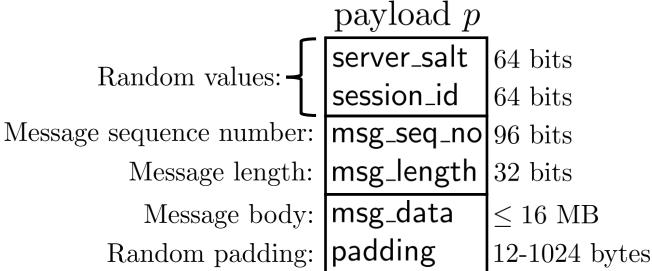
MAC – Message Authentication Code

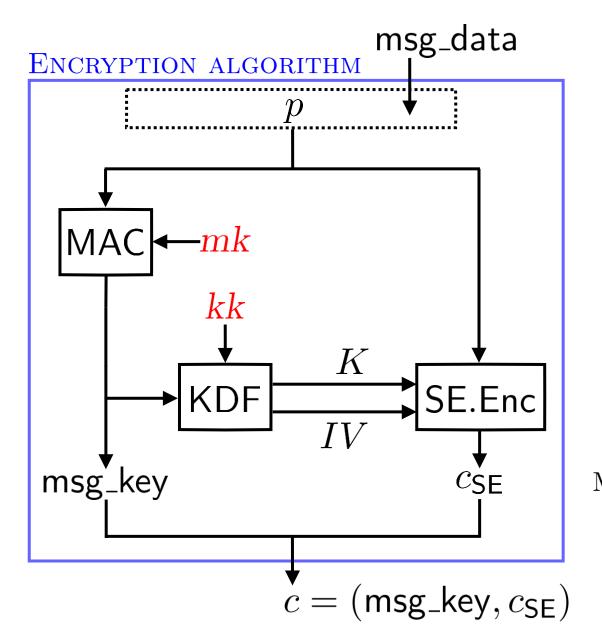
SE – Symmetric Encryption Scheme

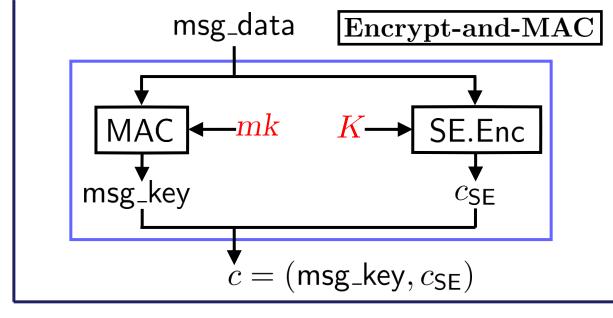




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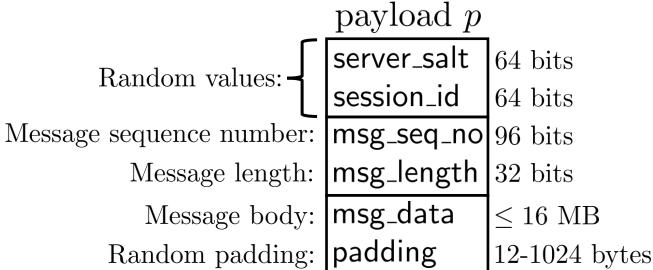


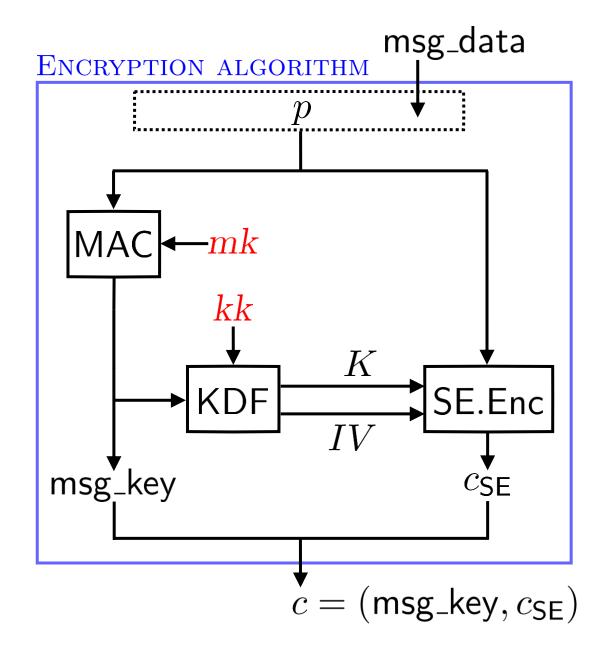


MAC – Message Authentication Code

SE – Symmetric Encryption Scheme

KDF – Key Derivation Function



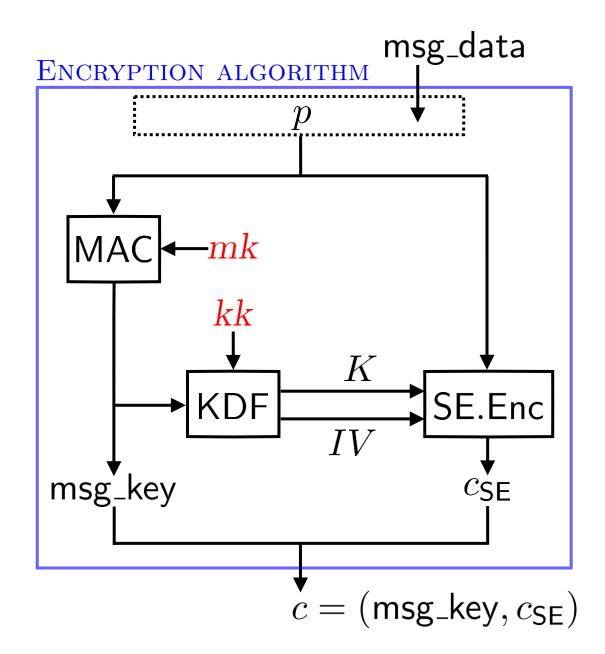


MTProto defines **ad-hoc** MAC and KDF schemes.

```
\frac{\mathsf{MAC}(\mathbf{mk}, p)}{\mathsf{msg\_key}} \leftarrow \mathsf{SHA-256}(\mathbf{mk} \| p)[64:192] Return \mathsf{msg\_key}
```

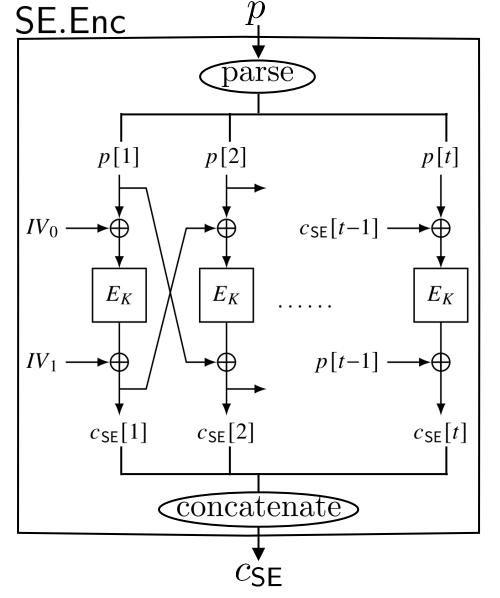
$$\frac{\mathsf{KDF}(kk, \mathsf{msg_key})}{(kk_0, kk_1) \leftarrow kk}$$
$$K \leftarrow \mathsf{SHA-256}(\mathsf{msg_key} || kk_0)$$
$$IV \leftarrow \mathsf{SHA-256}(kk_1 || \mathsf{msg_key})$$
$$\mathsf{Return} \ K, IV$$

Why invent new MAC and KDF schemes?

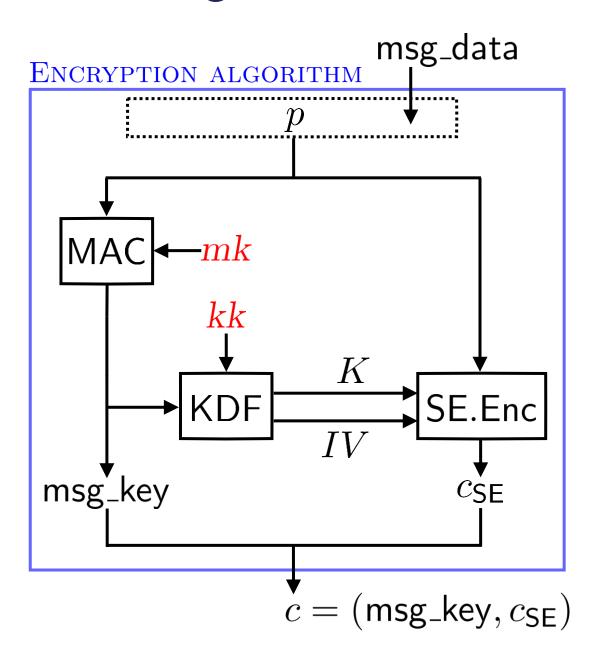


Infinite Garble Extension (IGE)

block cipher mode of operations

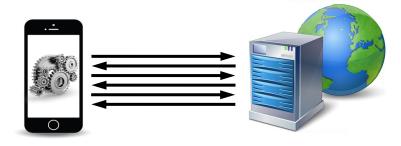


Not commonly used and not well studied.



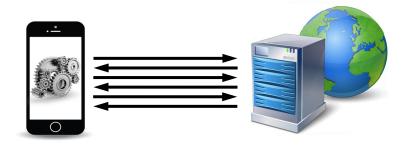
How are mk and kk derived?

MTProto uses Diffie-Hellman key exchange to agree on a raw shared secret g^{xy} .

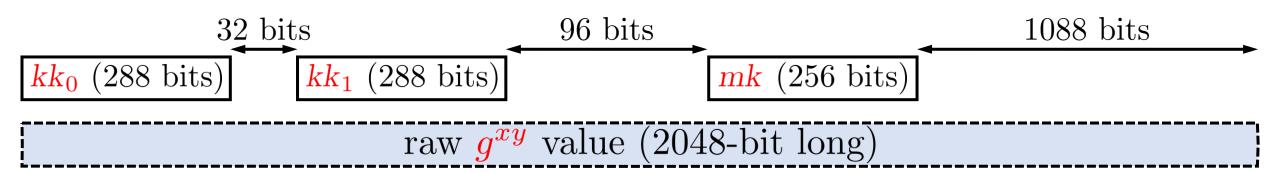


raw g^{xy} value (2048-bit long)

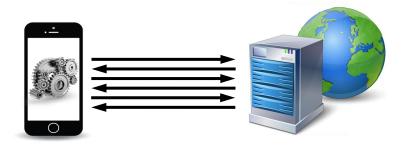
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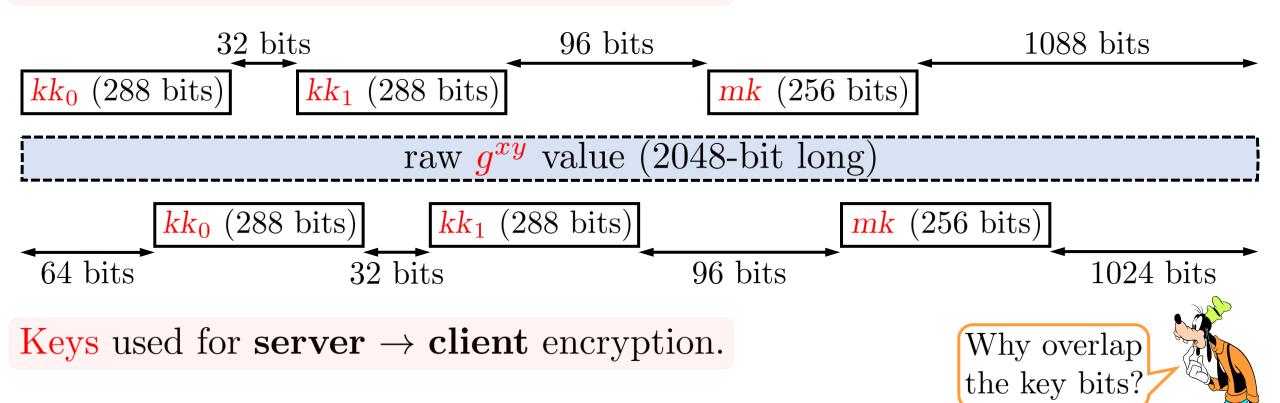
Keys used for client \rightarrow server encryption.



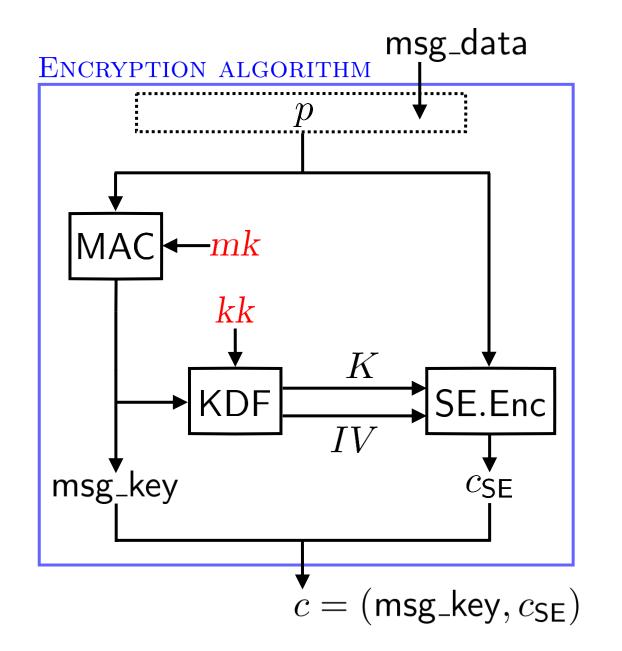
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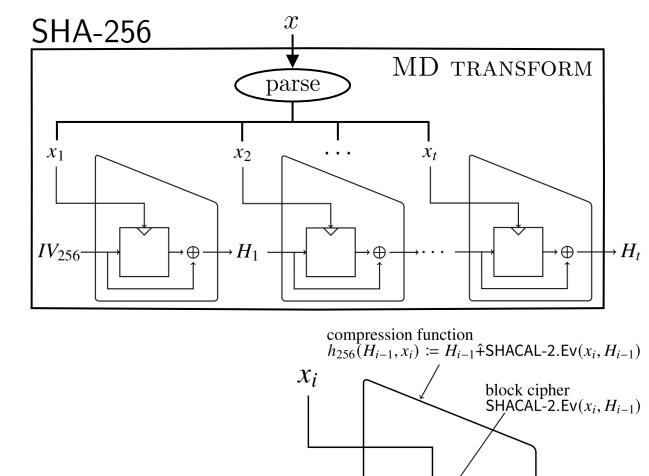


Keys used for **client** \rightarrow **server** encryption.



We proved a variant of MTProto 2.0 is secure. This comes with many caveats.



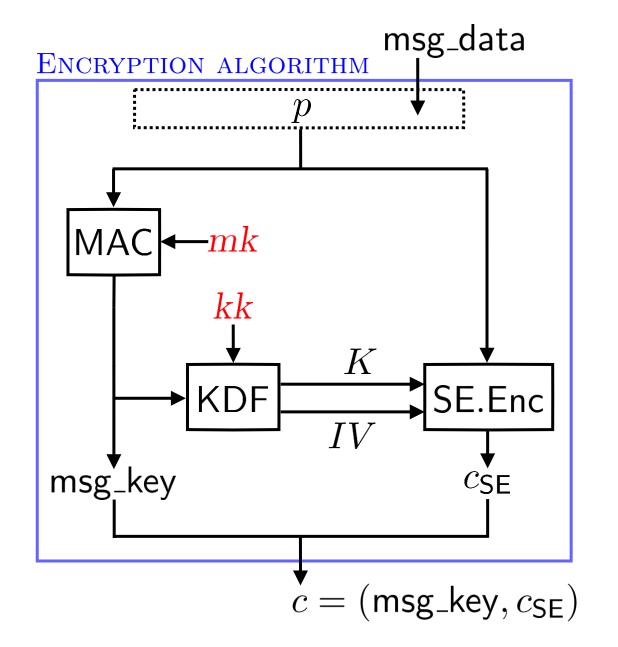


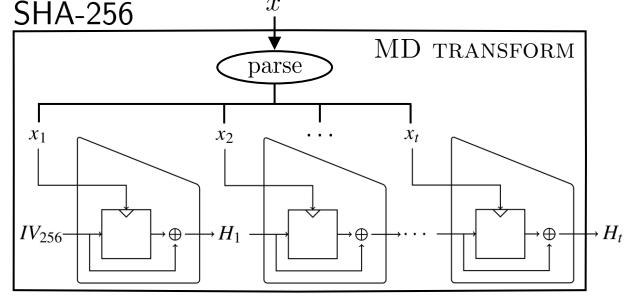
 H_{i-1}

 H_i

 \oplus

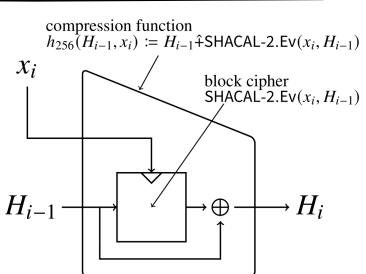
We proved a variant of MTProto 2.0 is secure. This comes with many caveats.





We rely on several assumptions about SHACAL-2.

Our novel SHACAL-2 assumptions need further study.



Four Attacks Against Telegram

April 16, 2021 We reported 4 vulnerabilities to **Telegram**.

April 22, 2021 Telegram confirmed the receipt of our e-mail.

June 08, 2021 Telegram acknowledged the reported behaviours.

July 16, 2021 Public disclosure (mutually agreed date).

2021 Telegram awarded bug bounty for attacks and analysis.

All vulnerabilities fixed as of

7.8.1 for **Android**

 $7.8.3 \text{ for } \mathbf{iOS}$

2.8.8 for **Desktop**





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Telegram informed us that they

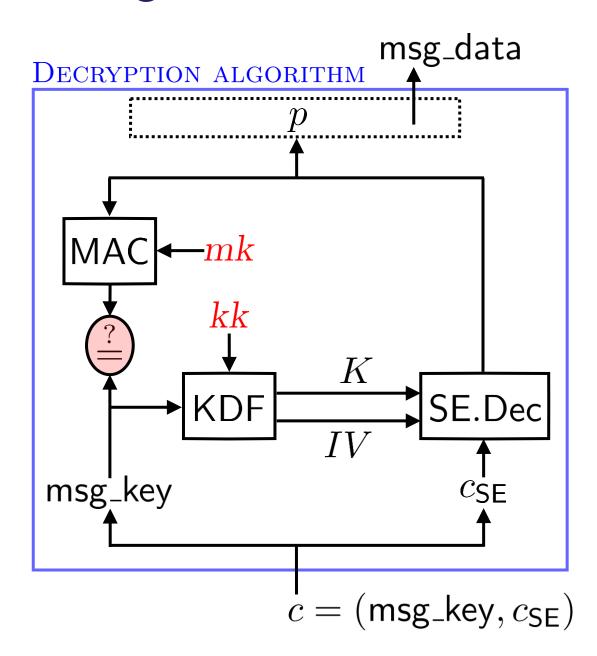
... do no security/bugfix releases except for post-release crash fixes.

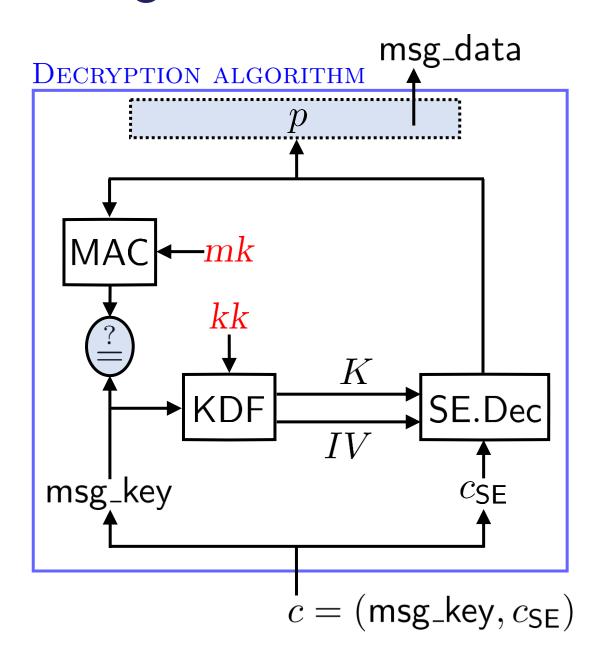
(could not commit to release dates for specific fixes)

(fixes were rolled out as part of regular updates)

... did not wish to issue security advisories at the time of patching.





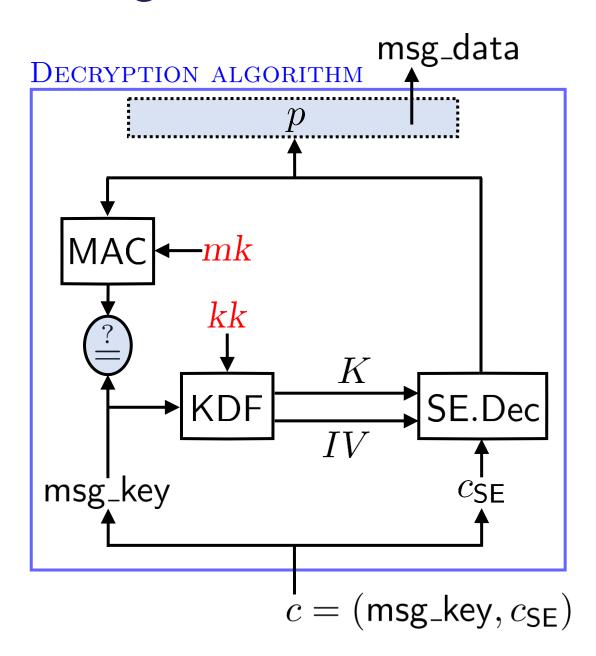


Telegram Desktop

```
if (msg\_length > 2^{24}) then

## MAC verification skipped
```

payload p server_salt 64 bits 64 bits 64 bits msg_seq_no 96 bits msg_length 32 bits msg_data padding ...



Telegram Desktop

if (msg_length > 2²⁴) then

MAC verification skipped

Introduces 3 microsecond difference.

Remote observer learns up to 8 bits.

payload p

server_salt 64 bits 64 bits 64 bits 64 bits 96 bits 96 bits 32 bits 65. msg_data padding ...

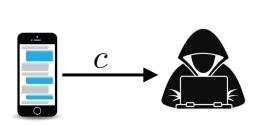
msg_length

8 upper bits 24 lower bits

We adapt the attack from:

Plaintext Recovery Attacks Against SSH

Martin R. Albrecht, Kenneth G. Paterson and Gaven J. Watson





Telegram Desktop

if (msg_length > 2²⁴) then

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	payload p	_
	server_salt	64 bits
	session_id	64 bits
	msg_seq_no	96 bits
	msg_length	32 bits
Э.	msg_data	
5.	padding	

msg_length

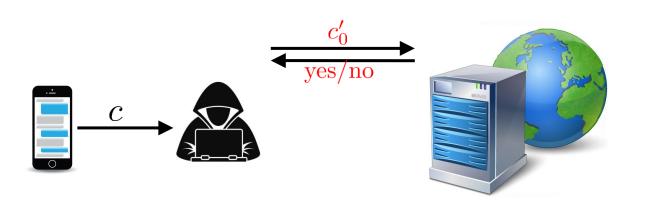
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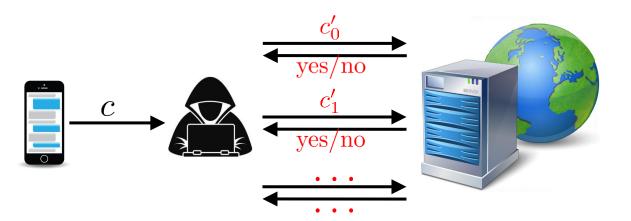
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"Are the upper 8 bits of msg_length (encrypted in c'_i) equal to 00000000?"

We recover (up to) 8 bits of msg_data per query!

Telegram Desktop

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msg_length

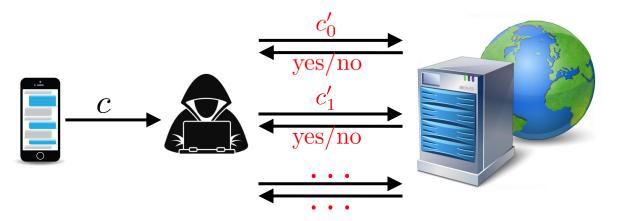
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msg_length

	<u> </u>
8 upper bits	24 lower bits

Three official clients checked p before MAC.

Telegram Desktop Telegram Android Telegram iOS

Each client did it in a different way.

Each client presented a timing side-channel.

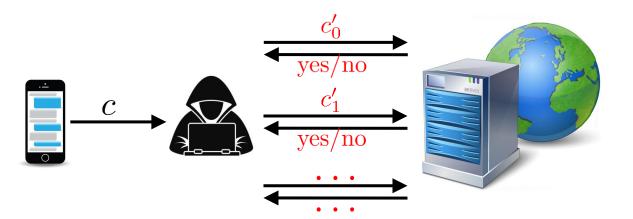
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msg_data ...
padding ...

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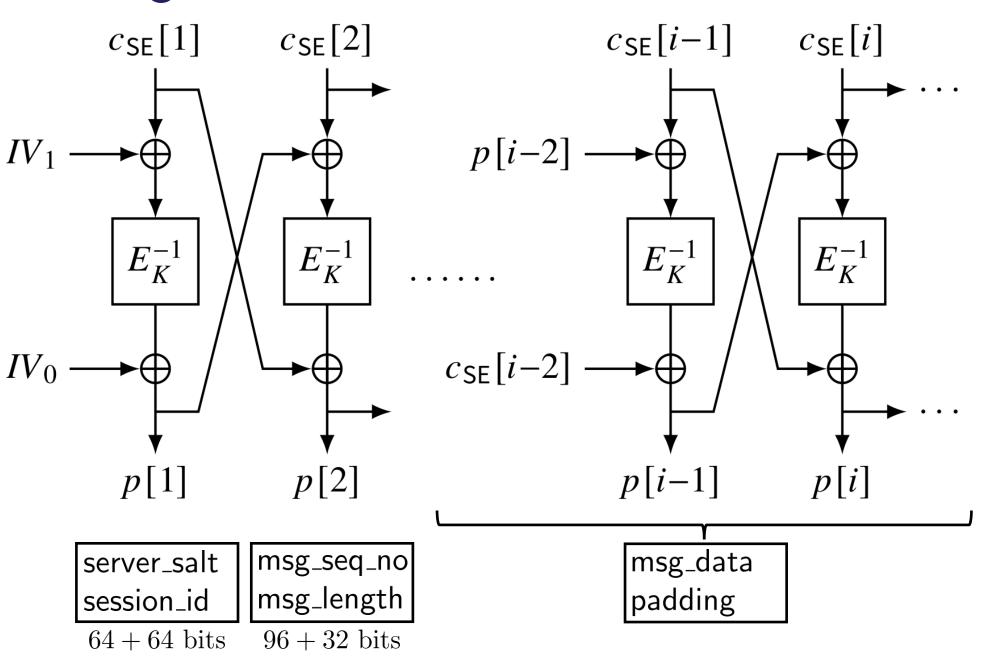
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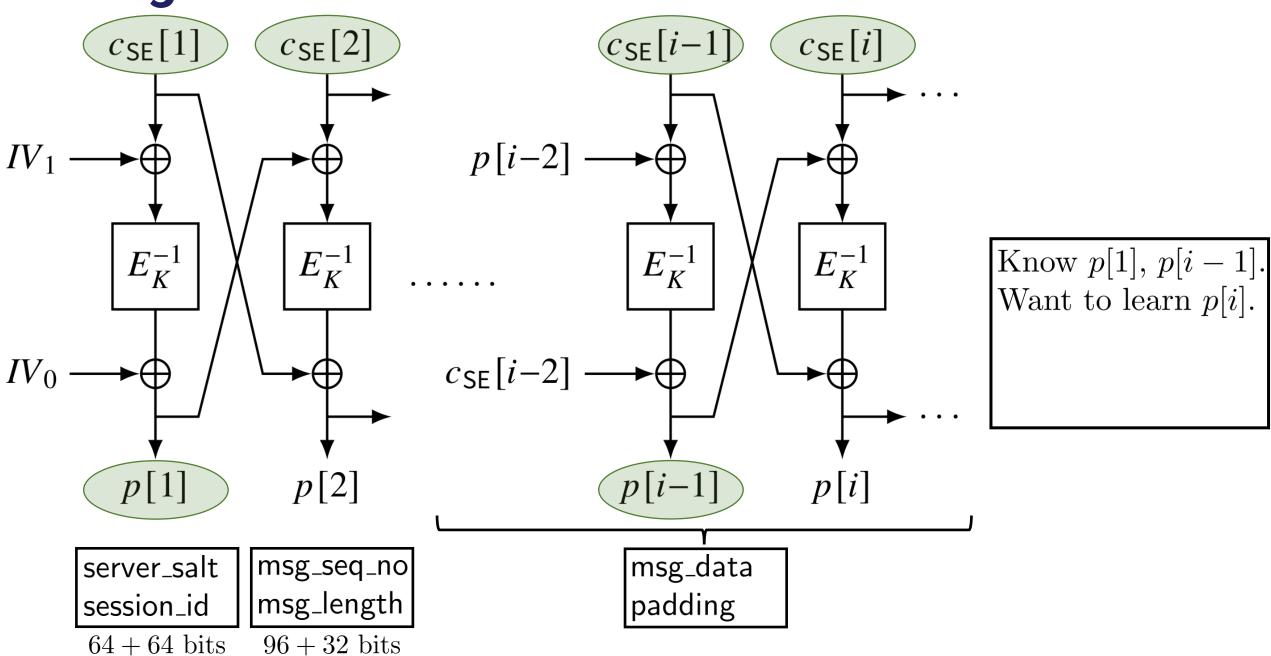
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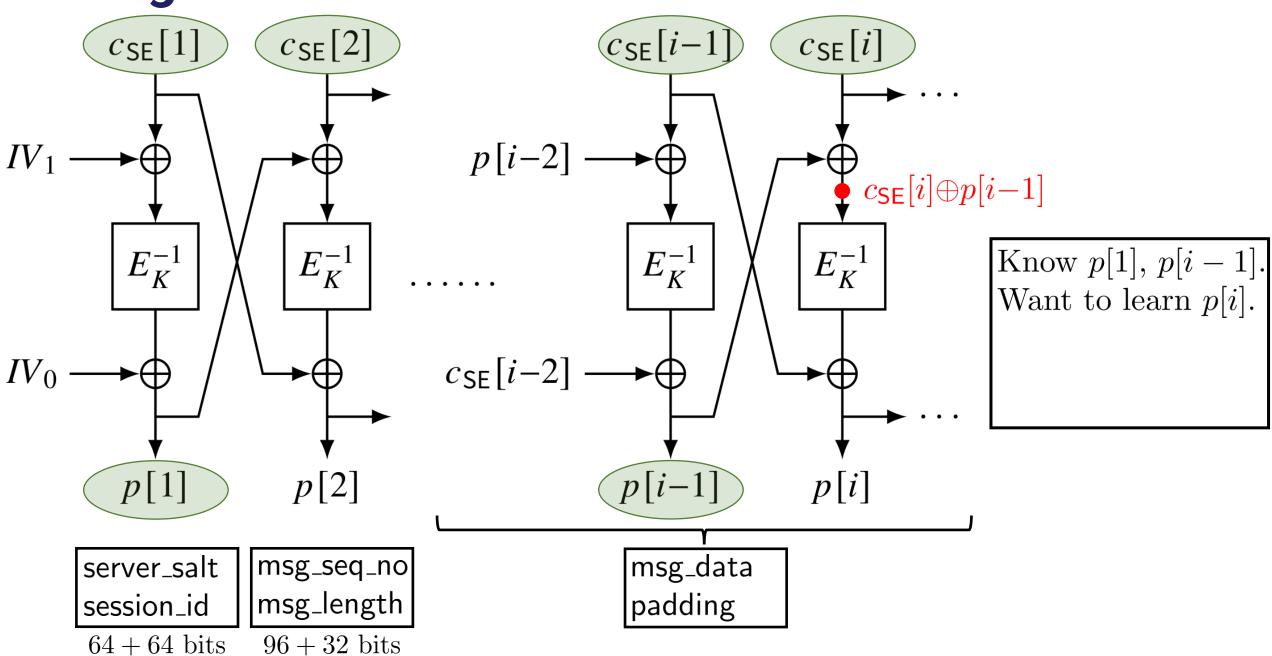
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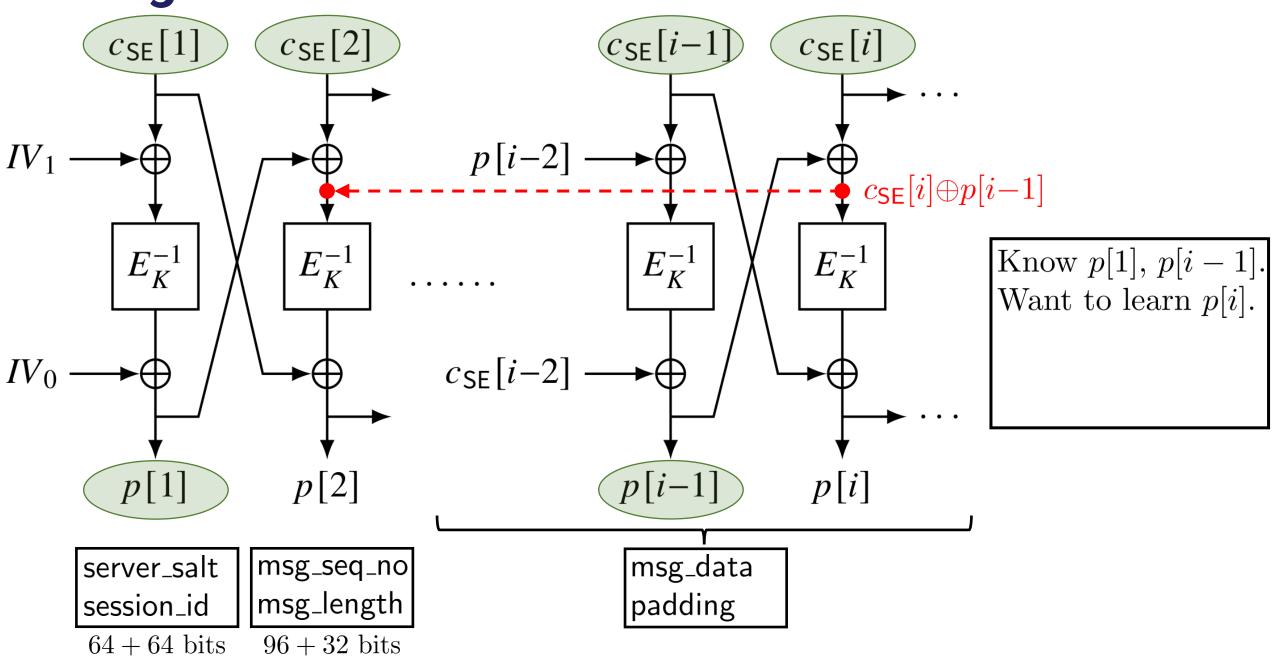
This highlights a brittle design.

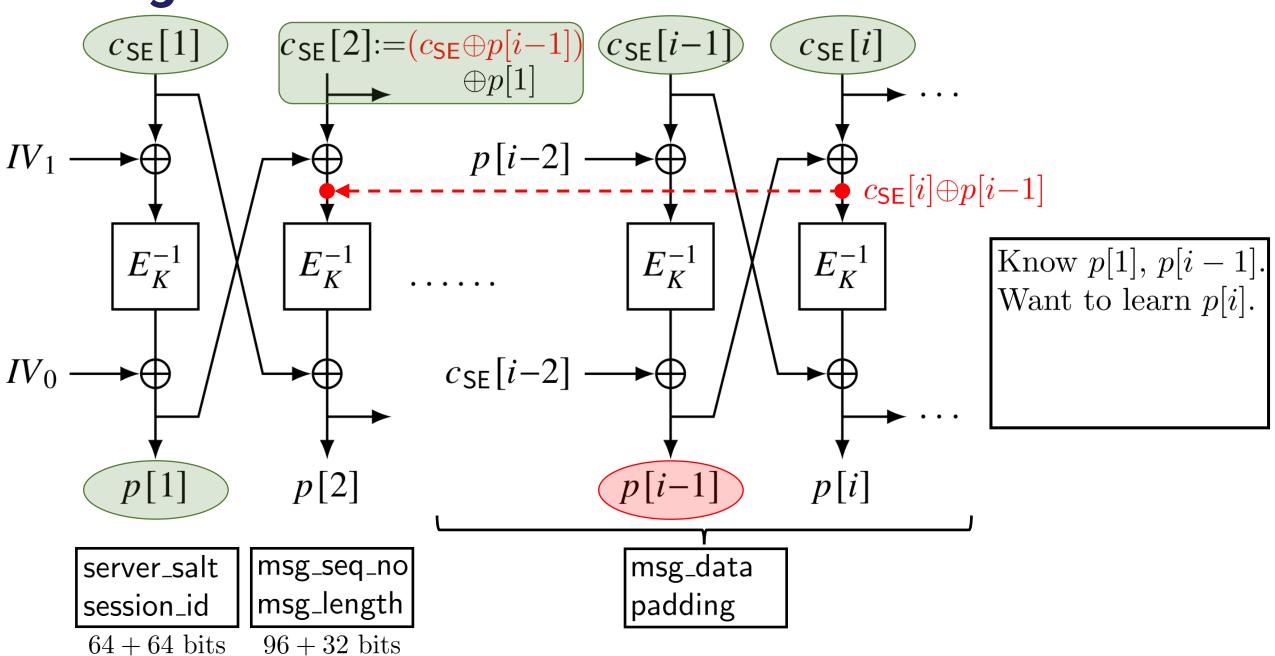
Encrypt-and-MAC requires to decrypt untrusted data. Would be safer to protect integrity of ciphertext.

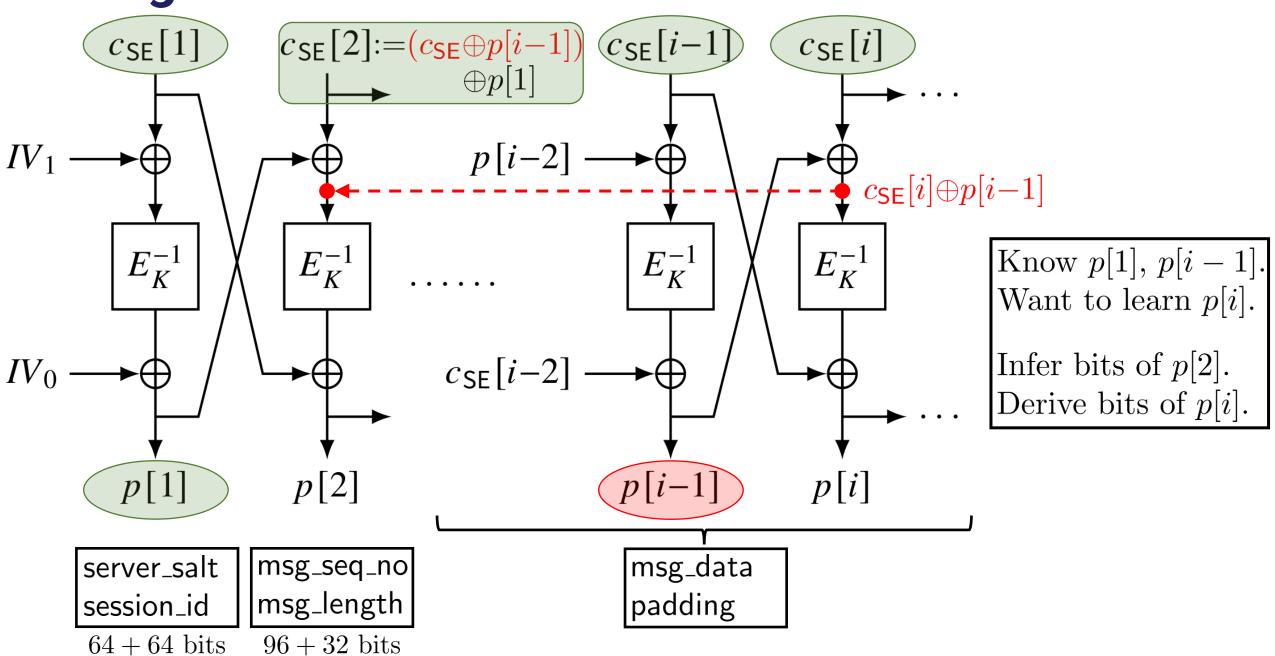


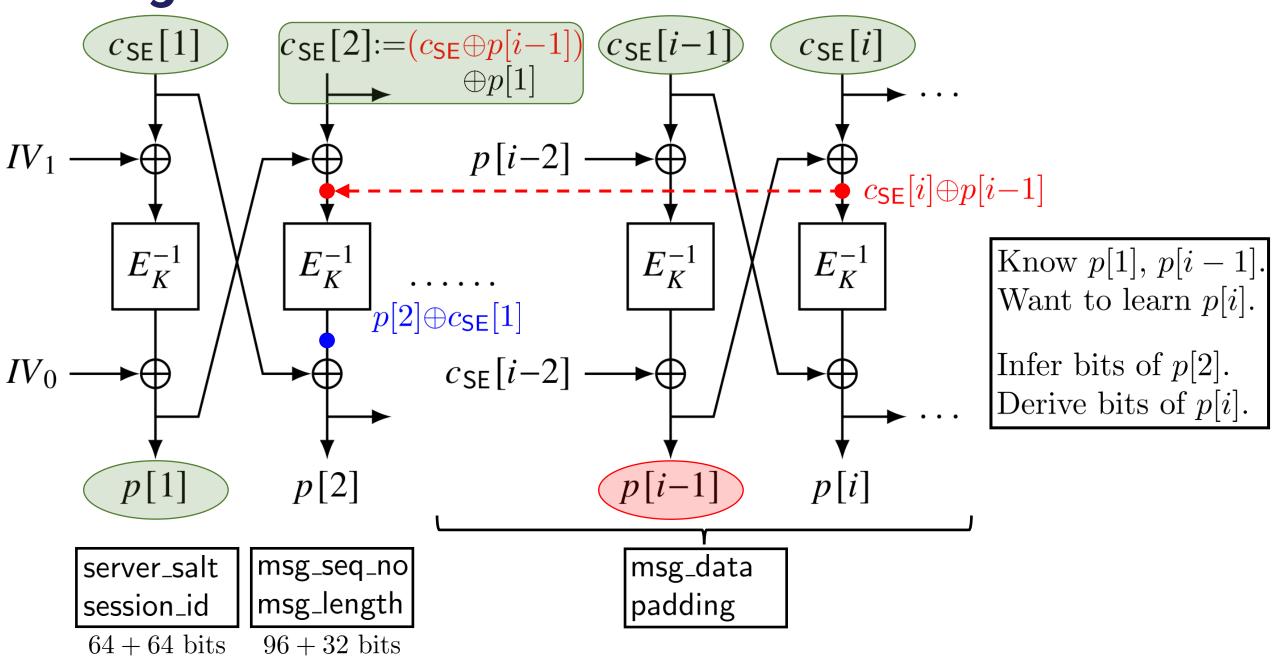


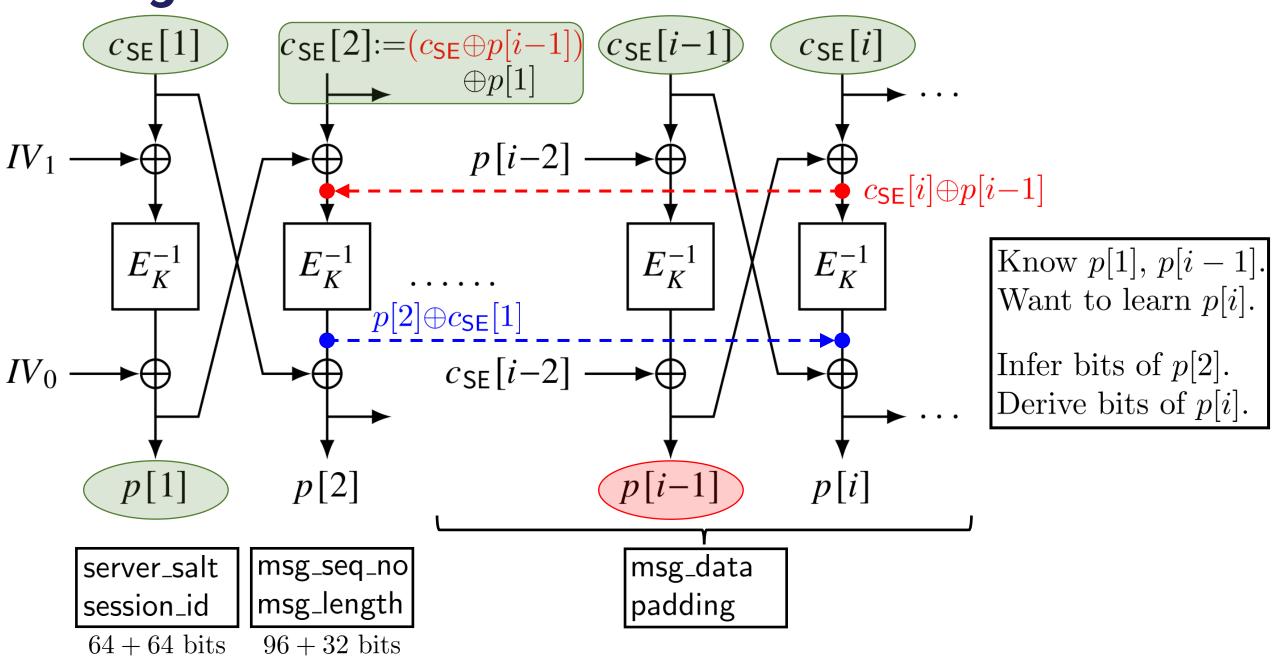


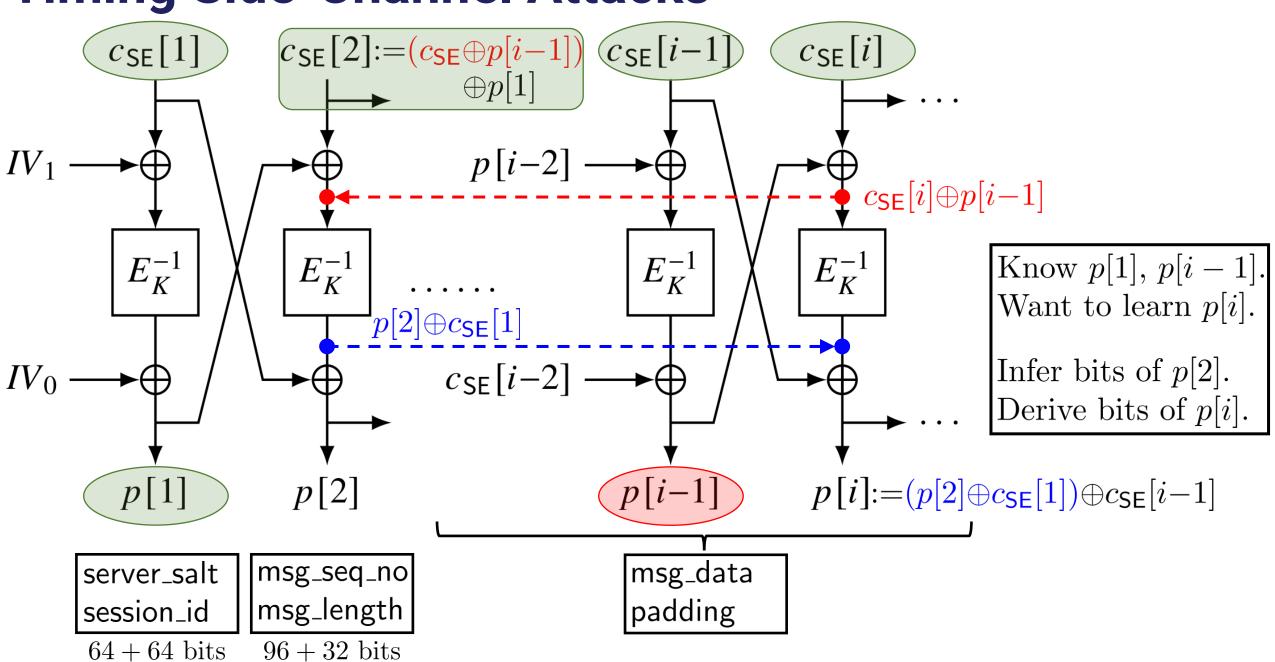


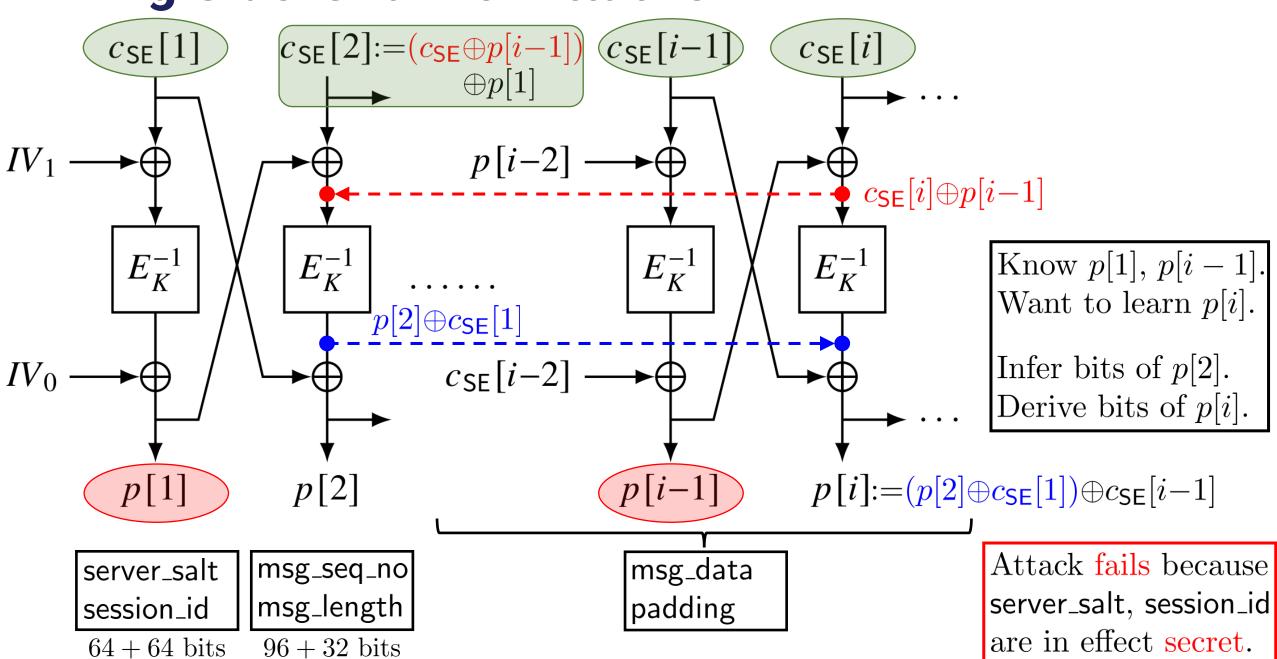




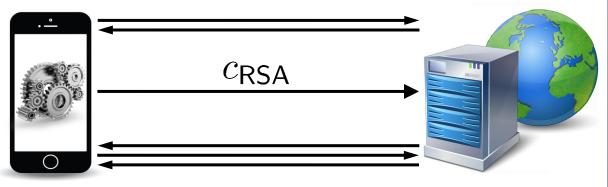






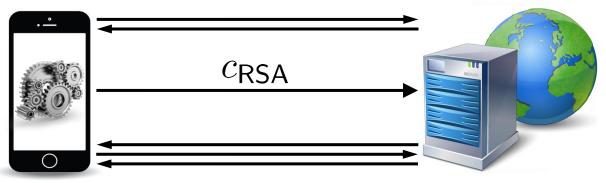


Telegram's key exchange



 $c_{\mathsf{RSA}} \leftarrow \mathsf{RSA}.\mathsf{Enc}(pk,n)$ $n = \mathsf{SHA-1}(\mathsf{data}) \| \mathsf{data} \| \mathsf{padding} \|$

Telegram's key exchange

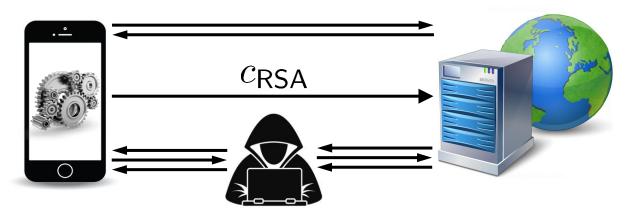


 $c_{\mathsf{RSA}} \leftarrow \mathsf{RSA}.\mathsf{Enc}(pk,n)$ $n = \mathsf{SHA}-1(\mathsf{data})\|\mathsf{data}\|\mathsf{padding}$ If data can be recovered:

Attacker learns server_salt immediately. Attacker learns session_id in $\approx 2^{64}$ queries.

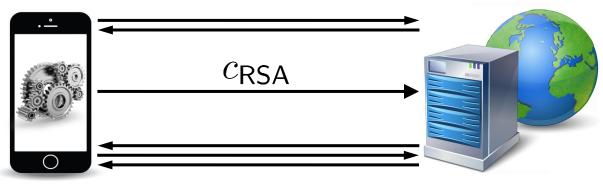
The attack against IGE now works!

If data can be recovered within 10 min:



Man-in-the-middle attack.

Telegram's key exchange



 $c_{\mathsf{RSA}} \leftarrow \mathsf{RSA}.\mathsf{Enc}(pk,n)$ $n = \mathsf{SHA-1}(\mathsf{data}) \| \mathsf{data} \| \mathsf{padding}$

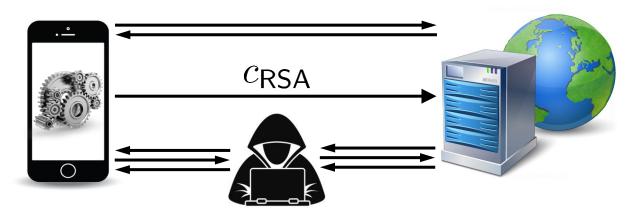
We recover data by solving noisy linear equations via lattice reduction.

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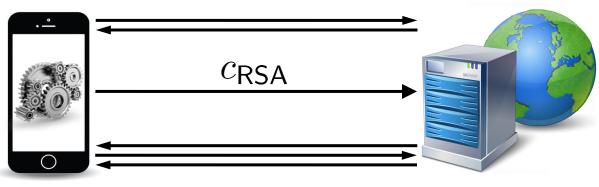
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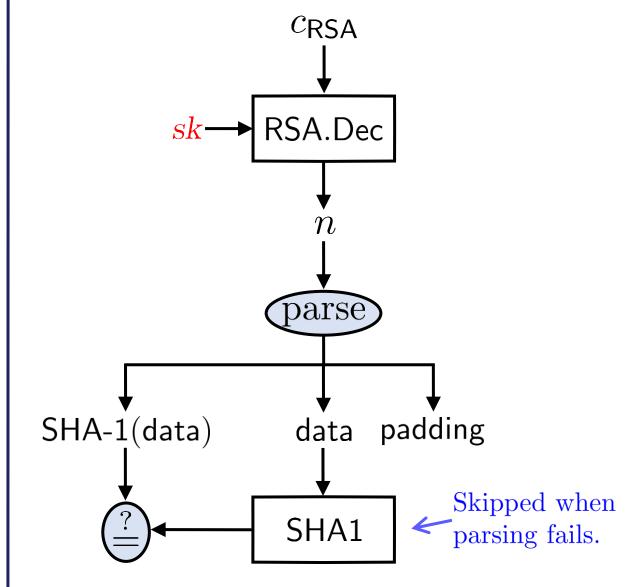
Telegram's key exchange



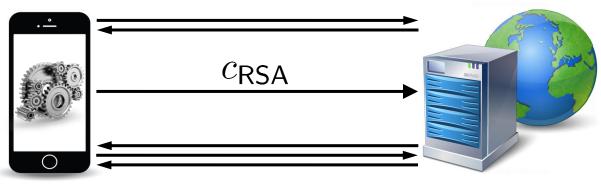
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Timing side-channel:



Telegram's key exchange

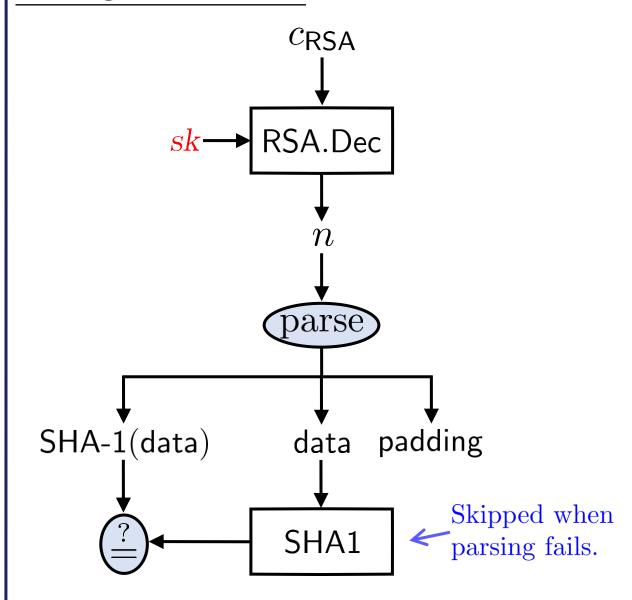


 $c_{\mathsf{RSA}} \leftarrow \mathsf{RSA}.\mathsf{Enc}(pk,n)$ $n = \mathsf{SHA-1}(\mathsf{data}) \| \mathsf{data} \| \mathsf{padding} \|$

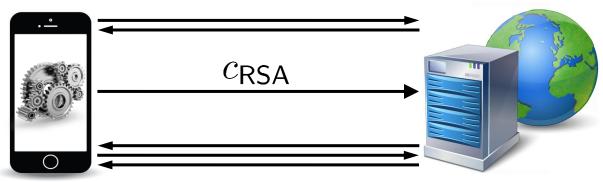
We recover data by solving noisy linear equations via lattice reduction.

Telegram uses textbook RSA scheme.
Why textbook RSA?

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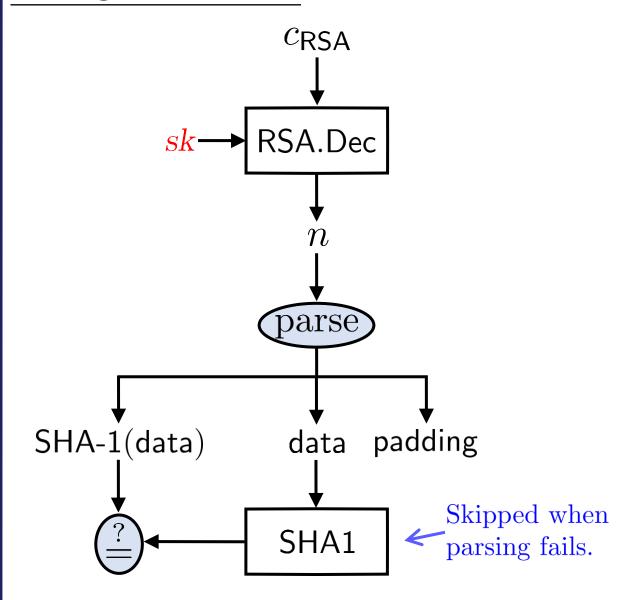
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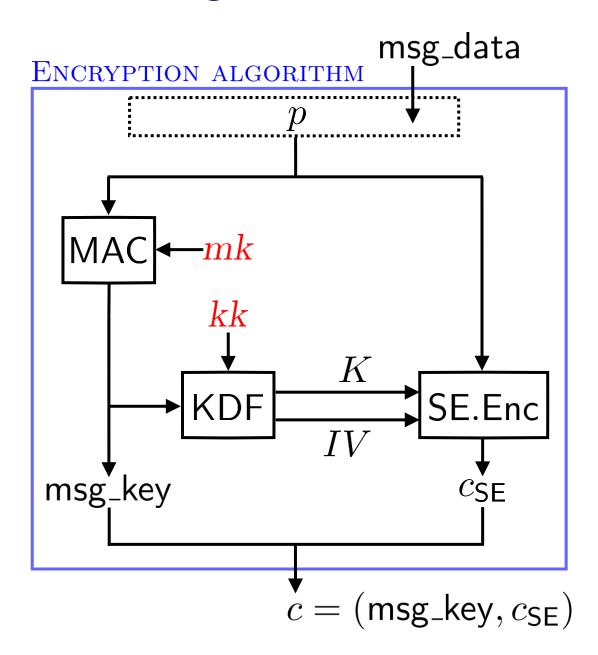
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Why textbook RSA?

Timings very small. Infeasible in practice. Caveat: Telegram's server code is secret.

Timing side-channel:

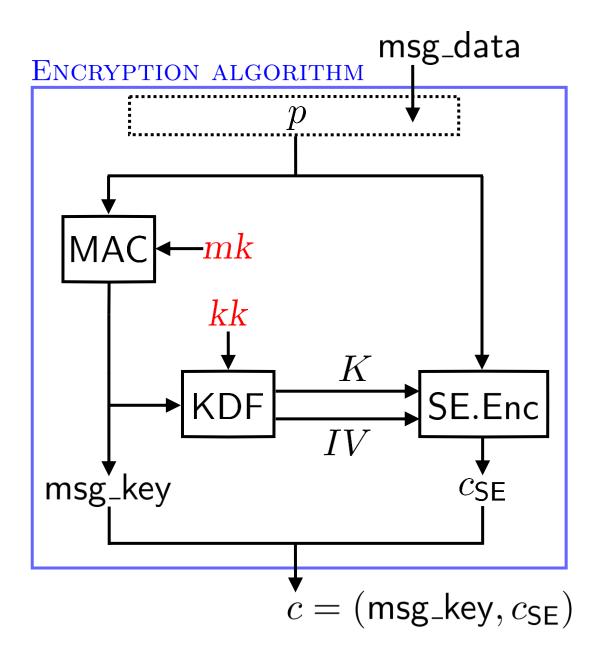




Theoretical attack.

MTProto requires message acknowledgements. If no acknowledgement, then payload is resent. Acknowledgements are encrypted.

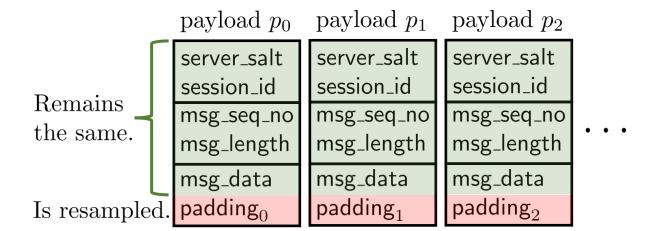
Our attack subverts this.

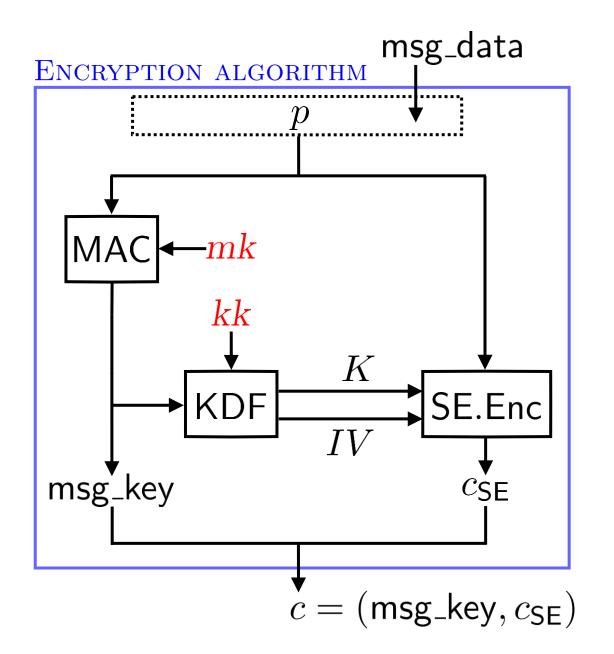


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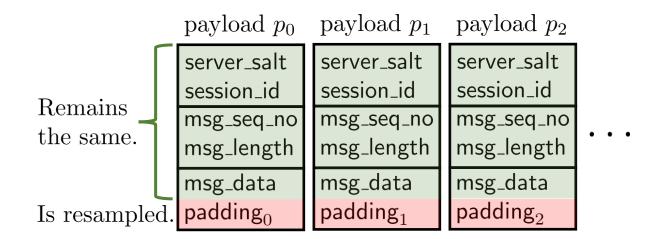




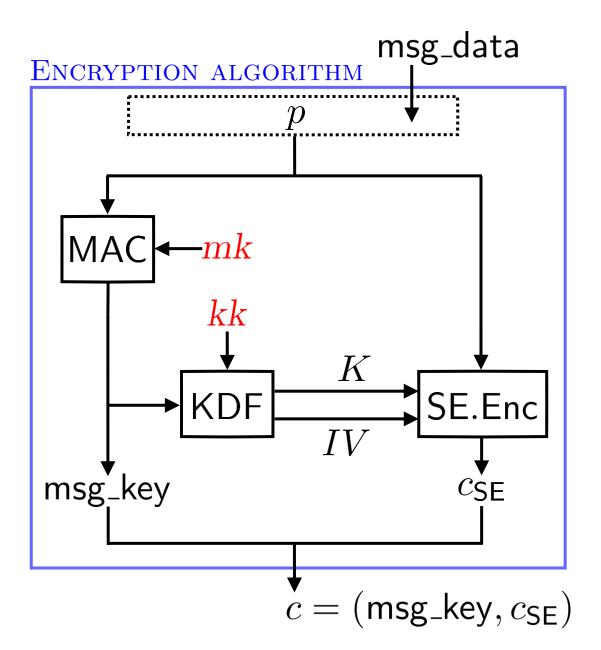
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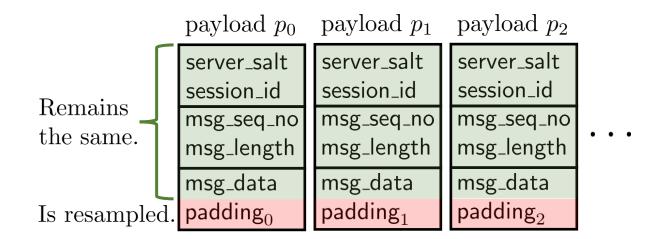
Observe
$$c = (\text{msg_key}, c_{SE}) \atop c^* = (\text{msg_key}^*, c_{SE}^*)$$
 with $c_{SE}[2] = c_{SE}^*[2]$?



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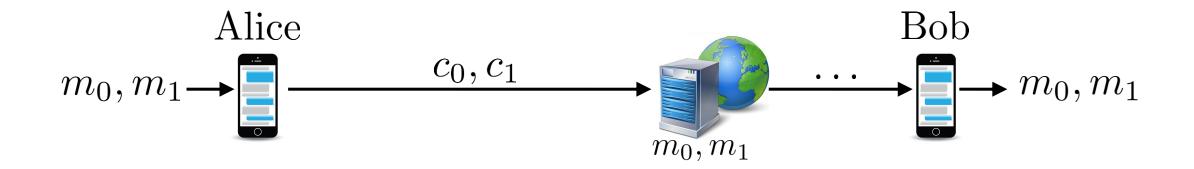


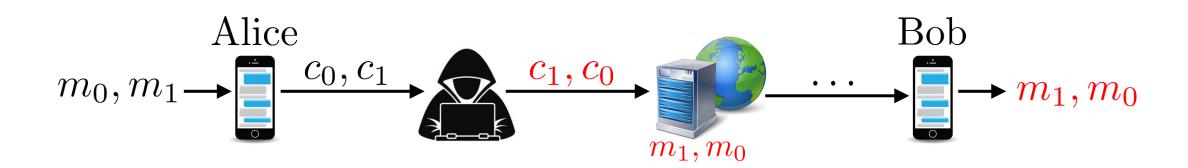
$$\textbf{Observe} \ \frac{c = (\mathsf{msg_key}, c_{SE})}{c^* = (\mathsf{msg_key}^*, c_{SE}^*)} \textbf{with} \ \frac{\mathsf{msg_key} = \mathsf{msg_key}^*}{c_{SE}[2] = c_{SE}^*[2]} \textbf{?}$$

Then c, c^* encrypt the same message!

No acknowledgement received between sending c, c^* .

Message Reordering Attack





Technically trivial. Easy to exploit.

Future Work

Large parts of **Telegram unstudied**:

SECRET CHATS
KEY EXCHANGE

..., multi-user security, forward secrecy, Telegram Passport, bot APIs, higher-level message processing, control messages, encrypted CDNs, cloud storage, ...



Thank you!

More information at https://mtpsym.github.io/